#### Towards Optimal Multi-Level Checkpointing

Anne Benoit<sup>1</sup>, Aurélien Cavelan<sup>1</sup>, Valentin Le Fèvre<sup>1</sup>, Yves Robert<sup>1,2</sup>, Hongyang Sun<sup>1</sup>

1. ENS Lyon & INRIA, France 2. University of Tennessee Knoxville, USA

hongyang.sun@ens-lyon.fr

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# Single-Level Checkpointing

Minimize expected execution overhead  $H(W) = \frac{\mathbb{E}(W)}{W} - 1$ 



Figure: periodic computing pattern

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• First-order approximation [Young/Daly]:

$$W_{\text{opt}} = \sqrt{\frac{2C}{\lambda}}$$
  
 $H_{\text{opt}} = \sqrt{2\lambda C} + \Theta(\lambda)$ 

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#### Scalability problem for large-scale platforms

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Two approaches:

• Independent checkpointing:



• Synchronized checkpointing:



Easy because pattern repeats (memoryless property)



Figure: with *N* level-1 checkpoints

• Exact solution: very complicated (which error type occurs first?), equal-length chunks, see [1]

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- First-order approximation:

$$H_{\mathrm{opt}} = \sqrt{2\lambda_1C_1} + \sqrt{2\lambda_2C_2} + \Theta(\lambda)$$

(obtained for some optimal pattern)

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• Choose optimal set of levels:

Level Overhead  
1, 2, 3 
$$\sqrt{2C_1\lambda_1} + \sqrt{2C_2\lambda_2} + \sqrt{2C_3\lambda_3}$$
  
1, 3  $\sqrt{2C_1\lambda_1} + \sqrt{2C_3(\lambda_2 + \lambda_3)}$   
2, 3  $\sqrt{2C_2(\lambda_1 + \lambda_2)} + \sqrt{2C_3\lambda_3}$   
3  $\sqrt{2C_3(\lambda_1 + \lambda_2 + \lambda_3)}$ 

## k Levels

#### Theorem

The optimal k-level pattern, under the first-order approximation, has equal-length chunks at all levels:

$$\begin{array}{l} \text{Optimal pattern length: } W^{opt} = \sqrt{\frac{\sum_{\ell=1}^{k} N_{\ell}^{opt} C_{\ell}}{\frac{1}{2} \sum_{\ell=1}^{k} \frac{\lambda_{\ell}}{N_{\ell}^{opt}}}} \\ \text{Optimal $\#$chkpts at level $\ell$: $N_{\ell}^{opt} = \sqrt{\frac{\lambda_{\ell}}{C_{\ell}} \cdot \frac{C_{k}}{\lambda_{k}}}, \quad \forall \ell = 1, \ldots, k \\ \text{Optimal pattern overhead: } H_{opt} = \sum_{\ell=1}^{k} \sqrt{2\lambda_{\ell}C_{\ell}} + \Theta(\lambda) \end{array}$$

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- Dynamic programming algorithm to choose set of levels
- Rounding for integer solution:  $n_{\ell}^{\text{opt}} = \frac{N_{\ell}^{\text{opt}}}{N_{\ell+1}^{\text{opt}}} = \sqrt{\frac{\lambda_{\ell}}{\lambda_{\ell+1}} \cdot \frac{C_{\ell+1}}{C_{\ell}}}$

#### Simulations

[	Set	Source	Level	1	2	3	4
Ī	(A)	Moody	C (s)	0.5	4.5	1051	-
		et al. [1]	MTBF (s)	5.00e6	5.56e5	2.50e6	-
[	(B)	Balaprakash	C (s)	10	20	20	100
		et al. [2]	MTBF (s)	3.60e4	7.20e4	1.44e5	7.20e5



[1] A. Moody, G. Bronevetsky, K. Mohror, and B. R. de Supinski. Design, modeling, and evaluation of a scalable multi-level checkpointing system. *Supercomputing*, 2010.

[2] P. Balaprakash, L. A. Bautista-Gomez, M.-S. Bouguerra, S. M. Wild, F. Cappello, and P. D. Hovland. Analysis of the tradeoffs between energy and run time for multilevel checkpointing. *PMBS*, 2014.

Explicit formulas for (almost) optimal multi-level checkpointing

$$H_{ ext{opt}} = \sum_{\ell=1}^k \sqrt{2\lambda_\ell C_\ell} + \Theta(\lambda)$$

Limitations:

- First-order approximation (accurate for 10,000s of nodes with MTBF in hours; beyond?)
- Independent errors (correlated failure?)